

This is the BirthdayBook specification, from Spivey [1].

$[NAME, DATE]$

The *BirthdayBook* schema defines the *state space* of the birthday book system.

$\textit{BirthdayBook}$
$known : \mathbb{P} NAME$ $birthday : NAME \mapsto DATE$
$known = \text{dom } birthday$

This *InitBirthdayBook* specifies the initial state of the birthday book system. It does not say explicitly that *birthday'* is empty, but that is implicit, because its domain is empty.

$\textit{InitBirthdayBook}$
$\textit{BirthdayBook}'$ $known' = \{\}$

Next we have several operation schemas to define the normal (non-error) behaviour of the system.

$\textit{AddBirthday}$
$\Delta \textit{BirthdayBook}$ $name? : NAME$ $date? : DATE$
$name? \notin known$ $birthday' = birthday \cup \{name? \mapsto date?\}$

$\textit{FindBirthday}$
$\exists \textit{BirthdayBook}$ $name? : NAME$ $date! : DATE$
$name? \in known$ $date! = birthday(name?)$

<i>Remind</i>
$\exists \text{BirthdayBook}$
$\text{today?} : \text{DATE}$
$\text{cards!} : \mathbb{P} \text{NAME}$
$\text{cards!} = \{n : \text{known} \mid \text{birthday}(n) = \text{today?}\}$

Now we strengthen the specification by adding error handling.

$\text{REPORT} ::= \text{ok} \mid \text{already\_known} \mid \text{not\_known}$

First we define auxiliary schemas that capture various success and error cases.

<i>Success</i>
$\text{result!} : \text{REPORT}$
$\text{result!} = \text{ok}$

<i>AlreadyKnown</i>
$\exists \text{BirthdayBook}$
$\text{name?} : \text{NAME}$
$\text{result!} : \text{REPORT}$
$\text{name?} \in \text{known}$
$\text{result!} = \text{already\_known}$

<i>NotKnown</i>
$\exists \text{BirthdayBook}$
$\text{name?} : \text{NAME}$
$\text{result!} : \text{REPORT}$
$\text{name?} \notin \text{known}$
$\text{result!} = \text{not\_known}$

Finally, we define robust versions of all the operations by specifying how errors are handled.

$R\text{AddBirthday} == (\text{AddBirthday} \wedge \text{Success}) \vee \text{AlreadyKnown}$   
 $R\text{FindBirthday} == (\text{FindBirthday} \wedge \text{Success}) \vee \text{NotKnown}$   
 $R\text{Remind} == \text{Remind} \wedge \text{Success}$

## References

- [1] J. Michael Spivey. *The Z Notation: A Reference Manual*. International Series in Computer Science. Prentice-Hall International (UK) Ltd, second edition, 1992.